

# Lecture 3

## Corner Polyhedron, Intersection Cuts, Maximal Lattice-Free Convex Sets

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# Mixed Integer Linear Programming

$$\begin{aligned} \min \quad & cx \\ \text{s.t.} \quad & Ax = b \\ & x_j \in \mathbb{Z} \quad \text{for } j = 1, \dots, p \\ & x_j \geq 0 \quad \text{for } j = 1, \dots, n. \end{aligned}$$

- First solve the LP relaxation. Basic optimal tableau:

$$x_i = \bar{b}_i - \sum_{j \in N} \bar{a}_{ij} x_j \quad \text{for } i \in B.$$

- If  $\bar{b}_i \notin \mathbb{Z}$  for some  $i \in B \cap \{1, \dots, p\}$ , add cutting planes:

For example Gomory 1963 Mixed Integer Cuts, Marchand and Wolsey 2001 MIR inequalities, Balas, Ceria and Cornuéjols 1993 lift-and-project cuts, are used in commercial codes.

# Corner Polyhedron [Gomory 1969]

Initial formulation:

$$\begin{aligned}x_i &= \bar{b}_i - \sum_{j \in N} \bar{a}_{ij} x_j && \text{for } i \in B. \\x_j &\in \mathbb{Z} && \text{for } j = 1, \dots, p \\x_j &\geq 0 && \text{for } j = 1, \dots, n.\end{aligned}$$

Corner formulation:

Drop the nonnegativity restriction on all the basic variables  $x_i$ ,  $i \in B$ .

**Note that** in this relaxation we can drop the constraints  $x_i = \bar{b}_i - \sum_{j \in N} \bar{a}_{ij} x_j$  for all  $i \in B \cap \{p+1, \dots, n\}$  since these variables  $x_i$  are continuous and only appear in one equation and no other constraint. Therefore from now on we assume that all basic variables in are integer, i.e.  $B \subseteq \{1, \dots, p\}$ .

# Corner Polyhedron

We assume  $B \subseteq \{1, \dots, p\}$ . The *corner formulation* is

$$\begin{aligned}x_i &= \bar{b}_i - \sum_{j \in N} \bar{a}_{ij} x_j && \text{for } i \in B \\x_i &\in \mathbb{Z} && \text{for } i = 1, \dots, p \\x_j &\geq 0 && \text{for } j \in N.\end{aligned}$$

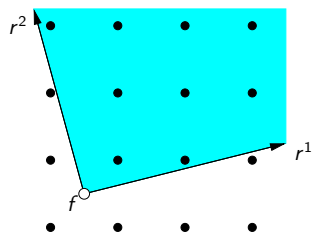
The convex hull of its feasible solutions is called the *corner polyhedron* relative to the basis  $B$  and it is denoted by  $\text{corner}(B)$ .

**Note** Any valid inequality for the corner polyhedron is valid for the initial formulation.

Let  $P(B)$  be the *linear relaxation* of the corner polyhedron.

$P(B)$  is a polyhedron whose vertices and extreme rays are simple to describe and this will be useful in generating valid inequalities for  $\text{corner}(B)$ .

# Corner Polyhedron Example



Feasible set  $\left\{ \begin{pmatrix} x_3 \\ x_4 \end{pmatrix} \in \mathbb{Z}^2 : \right.$

$$\begin{pmatrix} x_3 \\ x_4 \end{pmatrix} = f + r^1 x_1 + r^2 x_2$$

where  $x_1 \geq 0, x_2 \geq 0$

Restricted to the  $(x_3, x_4)$ -space,  $P(B)$  is the blue region.  
The feasible solutions are the integer points in the blue region,  
and  $\text{corner}(B)$  is the convex hull of these points.

## The Polyhedron $P(B)$

$P(B)$  has a unique vertex  $\bar{x}$  where  $\bar{x}_i = \bar{b}_i, i \in B, x_j = 0, j \in N$ .

The recession cone of  $P(B)$  is

$$\begin{aligned} x_i &= -\sum_{j \in N} \bar{a}_{ij} x_j && \text{for } i \in B \\ x_j &\geq 0 && \text{for } j \in N. \end{aligned}$$

Since the projection of this cone onto  $\mathbb{R}^N$  is defined by the inequalities  $x_j \geq 0, j \in N$ , its extreme rays are the vectors satisfying at equality all but one nonnegativity constraints. Thus there are  $|N|$  extreme rays,  $\bar{r}^j$  for  $j \in N$ , defined by

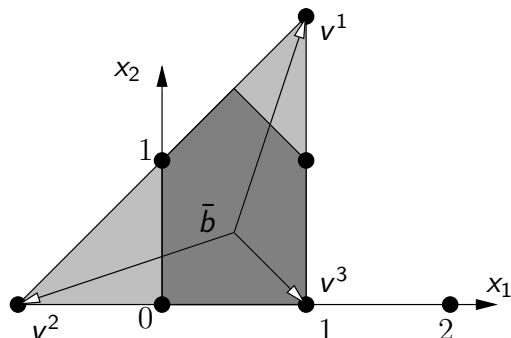
$$\bar{r}_h^j = \begin{cases} -\bar{a}_{hj} & \text{if } h \in B, \\ 1 & \text{if } h = j, \\ 0 & \text{if } h \in N \setminus \{j\}. \end{cases}$$

**Remark** The vectors  $\bar{r}^j, j \in N$  are linearly independent. Hence  $P(B)$  is an  $|N|$ -dimensional polyhedron whose affine hull is defined by the equations  $x_i = \bar{b}_i - \sum_{j \in N} \bar{a}_{ij} x_j$  for  $i \in B$ .

# Corner Polyhedron Example

Consider the pure integer program

$$\begin{aligned} \max \quad & \frac{1}{2}x_2 + x_3 \\ & x_1 + x_2 + x_3 \leq 2 \\ & x_1 - \frac{1}{2}x_3 \geq 0 \\ & x_2 - \frac{1}{2}x_3 \geq 0 \\ & x_1 + \frac{1}{2}x_3 \leq 1 \\ & -x_1 + x_2 + x_3 \leq 1 \\ & x_1, x_2, x_3 \in \mathbb{Z} \\ & x_1, x_2, x_3 \geq 0. \end{aligned}$$

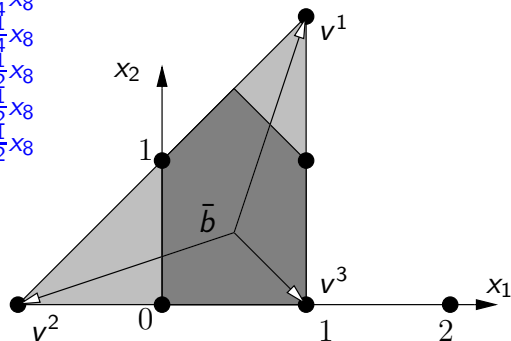


This problem has 4 feasible solutions  $(0, 0, 0)$ ,  $(1, 0, 0)$ ,  $(0, 1, 0)$  and  $(1, 1, 0)$ , all satisfying  $x_3 = 0$ . The intersection of the 5 inequalities in the formulation with the plane  $x_3 = 0$  is the darker region in the figure.

## Corner Polyhedron Example

We first write the problem in standard form by introducing continuous slack or surplus variables  $x_4, \dots, x_8$ . Solving the LP relaxation, we get

$$\begin{aligned}x_1 &= \frac{1}{2} + \frac{1}{4}x_6 - \frac{3}{4}x_7 + \frac{1}{4}x_8 \\x_2 &= \frac{1}{2} + \frac{3}{4}x_6 - \frac{1}{4}x_7 - \frac{1}{4}x_8 \\x_3 &= 1 - \frac{1}{2}x_6 - \frac{1}{2}x_7 - \frac{1}{2}x_8 \\x_4 &= 0 - \frac{1}{2}x_6 + \frac{3}{2}x_7 + \frac{1}{2}x_8 \\x_5 &= 0 + \frac{1}{2}x_6 - \frac{1}{2}x_7 + \frac{1}{2}x_8\end{aligned}$$



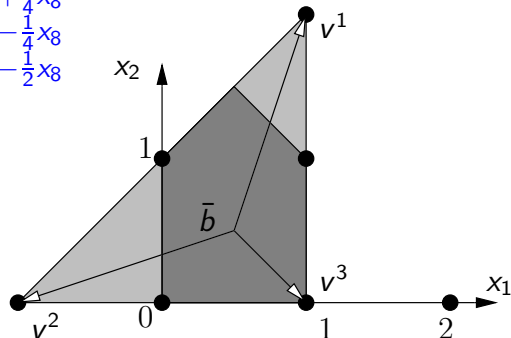
The optimal basic solution is  $x_1 = x_2 = \frac{1}{2}$ ,  $x_3 = 1$ ,  
 $x_4 = \dots = x_8 = 0$ .



# Corner Polyhedron Example

Relaxing the nonnegativity of the basic variables and dropping the two constraints relative to the continuous basic variables  $x_4$  and  $x_5$ , we obtain the **corner formulation**:

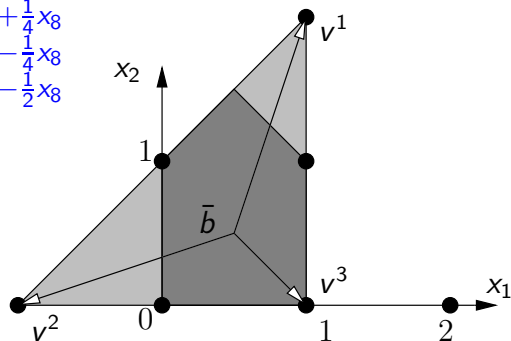
$$\begin{array}{rcll} x_1 & = & \frac{1}{2} & + \frac{1}{4}x_6 & - \frac{3}{4}x_7 & + \frac{1}{4}x_8 \\ x_2 & = & \frac{1}{2} & + \frac{3}{4}x_6 & - \frac{1}{4}x_7 & - \frac{1}{4}x_8 \\ x_3 & = & 1 & - \frac{1}{2}x_6 & - \frac{1}{2}x_7 & - \frac{1}{2}x_8 \\ x_1, x_2, x_3 & \in & & & \mathbb{Z} & \\ x_6, x_7, x_8 & \geq & & & 0. & \end{array}$$



## Corner Polyhedron Example

Let  $P(B)$  be the linear relaxation of corner formulation. The projection of  $P(B)$  in the space of original variables  $x_1, x_2, x_3$  is a polyhedron with unique vertex  $\bar{b} = (\frac{1}{2}, \frac{1}{2}, 1)$ . The extreme rays of its recession cone are  $v^1 = (\frac{1}{2}, \frac{3}{2}, -1)$ ,  $v^2 = (-\frac{3}{2}, -\frac{1}{2}, -1)$  and  $v^3 = (\frac{1}{2}, -\frac{1}{2}, -1)$ .

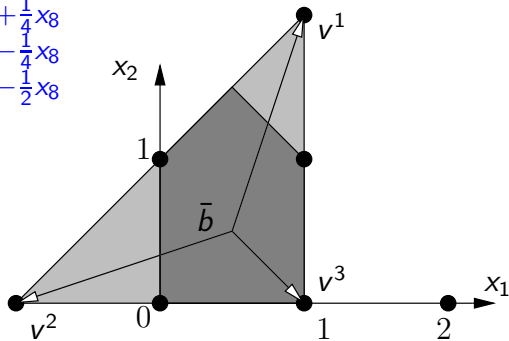
$$\begin{array}{rcll} x_1 & = & \frac{1}{2} & + \frac{1}{4}x_6 & - \frac{3}{4}x_7 & + \frac{1}{4}x_8 \\ x_2 & = & \frac{1}{2} & + \frac{3}{4}x_6 & - \frac{1}{4}x_7 & - \frac{1}{4}x_8 \\ x_3 & = & 1 & - \frac{1}{2}x_6 & - \frac{1}{2}x_7 & - \frac{1}{2}x_8 \\ x_6, x_7, x_8 & & & \geq & & 0. \end{array}$$



## Corner Polyhedron Example

In the figure, the shaded region (both light and dark) is the intersection of  $P(B)$  with the plane  $x_3 = 0$ .

$$\begin{array}{rcll} x_1 & = & \frac{1}{2} & + \frac{1}{4}x_6 & - \frac{3}{4}x_7 & + \frac{1}{4}x_8 \\ x_2 & = & \frac{1}{2} & + \frac{3}{4}x_6 & - \frac{1}{4}x_7 & - \frac{1}{4}x_8 \\ x_3 & = & 1 & - \frac{1}{2}x_6 & - \frac{1}{2}x_7 & - \frac{1}{2}x_8 \\ x_6, x_7, x_8 & & & \geq & & 0. \end{array}$$



Let  $P$  be defined by the inequalities of the initial formulation that are satisfied at equality by the point  $\bar{b} = (\frac{1}{2}, \frac{1}{2}, 1)$ . The intersection of  $P$  with the plane  $x_3 = 0$  is the dark shaded region.

## Intersection Cuts [Balas 1971]

Consider a closed convex set  $C \subseteq \mathbb{R}^n$  such that the interior of  $C$  contains  $\bar{x}$  but no point in  $\mathbb{Z}^p \times \mathbb{R}^{n-p}$ .

For each of the  $|N|$  extreme rays of  $P(B)$ , define

$$\alpha_j = \max\{\alpha \geq 0 : \bar{x} + \alpha \bar{r}^j \in C\}.$$

Since  $\bar{x}$  is in the interior of  $C$ ,  $\alpha_j > 0$ .

When the half-line  $\{\bar{x} + \alpha \bar{r}^j : \alpha \geq 0\}$  intersects the boundary of  $C$ , then  $\alpha_j$  is finite, and the point  $\bar{x} + \alpha_j \bar{r}^j$  belongs to the boundary of  $C$ .

When  $\bar{r}^j$  belongs the recession cone of  $C$ , we have  $\alpha_j = +\infty$ .

Define  $\frac{1}{+\infty} = 0$ .

The inequality

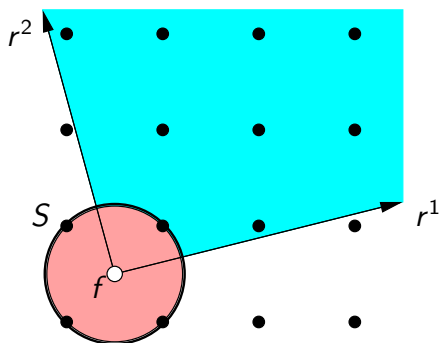
$$\sum_{j \in N} \frac{x_j}{\alpha_j} \geq 1$$

is the *intersection cut* defined by  $C$ .

# Intersection Cuts

Assume  $f \notin \mathbb{Z}^2$ .

Want to cut off the basic solution  $\begin{pmatrix} x_3 \\ x_4 \end{pmatrix} = f, x_1 = 0, x_2 = 0$ .



Feasible set  $\left\{ \begin{pmatrix} x_3 \\ x_4 \end{pmatrix} \in \mathbb{Z}^2 : \right.$

$$\begin{pmatrix} x_3 \\ x_4 \end{pmatrix} = f + r^1 x_1 + r^2 x_2$$

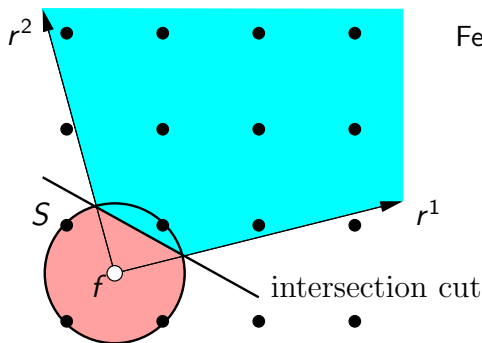
where  $x_1 \geq 0, x_2 \geq 0$

Any convex set  $S$  with  $f \in \text{int}(S)$  and no integer point in  $\text{int}(S)$ .

# Intersection Cuts

Assume  $f \notin \mathbb{Z}^2$ .

Want to cut off the basic solution  $\begin{pmatrix} x_3 \\ x_4 \end{pmatrix} = f, x_1 = 0, x_2 = 0$ .



Feasible set  $\left\{ \begin{pmatrix} x_3 \\ x_4 \end{pmatrix} \in \mathbb{Z}^2 : \right.$

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where  $x_1 \geq 0, x_2 \geq 0$

Any convex set  $S$  with  $f \in \text{int}(S)$  and no integer point in  $\text{int}(S)$ .

*Intersection cut* is obtained by intersecting the rays with the boundary of  $S$ :  $\alpha_1 = \frac{1}{4}, \alpha_2 = \frac{1}{4}$ . Thus  $4x_1 + 4x_2 \geq 1$ .

# Intersection Cuts

The *corner formulation* introduced by **Gomory** is

$$\begin{aligned}x_i &= \bar{b}_i - \sum_{j \in N} \bar{a}_{ij} x_j && \text{for } i \in B \\x_i &\in \mathbb{Z} && \text{for } i = 1, \dots, p \\x_j &\geq 0 && \text{for } j \in N.\end{aligned}$$

Basic solution  $\bar{x}$  where  $\bar{x}_i = \bar{b}_i, i \in B, x_j = 0, j \in N$ .

Assume  $B \subseteq \{1, \dots, p\}$  and  $\bar{b} \notin \mathbb{Z}^{|B|}$ .

**THEOREM** Let  $C \subset \mathbb{R}^n$  be a closed convex set whose interior contains the point  $\bar{x}$  but no point in  $\mathbb{Z}^p \times \mathbb{R}^{n-p}$ . The intersection cut defined by  $C$  is a valid inequality for  $\text{corner}(B)$ .

**PROOF** See course notes

# Intersection Cuts $\equiv$ Corner Polyhedron

We just stated that intersection cuts are valid for  $\text{corner}(B)$ .

The following theorem provides a converse statement.

Inequalities  $\sum_{j \in N} \gamma_j x_j \geq \gamma_0$  with  $\gamma_j \geq 0, j \in N$  and  $\gamma_0 \leq 0$  are implied by the nonnegativity constraints  $x_j \geq 0, j \in N$  and will be called *trivial*.

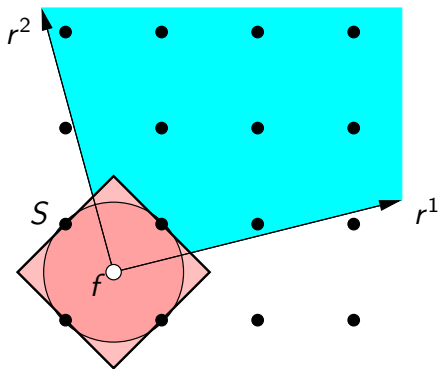
Every nontrivial valid inequality for  $\text{corner}(B)$  can be written in the form  $\sum_{j \in N} \gamma_j x_j \geq 1$  with  $\gamma_j \geq 0, j \in N$ .

**THEOREM** Every nontrivial facet defining inequality for  $\text{corner}(B)$  is an intersection cut.

**PROOF** See course notes



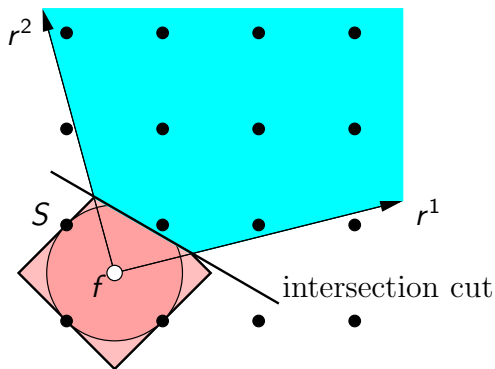
## A Better Intersection Cut for our Example



Bigger convex set:

Square  $S$  such that  $f \in \text{int}(S)$  with no integral point in  $\text{int}(S)$ .

## A Better Intersection Cut for our Example



Bigger convex set:

Square  $S$  such that  $f \in \text{int}(S)$  with no integral point in  $\text{int}(S)$ .

Better cut:  $\alpha_1 = \frac{1}{3}$ ,  $\alpha_2 = \frac{1}{3}$ . Thus  $3x_1 + 3x_2 \geq 1$ .

# Intersection Cuts

- If  $C_1 \subset C_2$  are two closed convex sets whose interiors contain  $\bar{x}$  but no point of  $\mathbb{Z}^p \times \mathbb{R}^{n-p}$ , then the intersection cut relative to  $C_2$  dominates the intersection cut relative to  $C_1$  for all  $x \in \mathbb{R}^n$  such that  $x_j \geq 0, j \in N$ .
- Any closed convex set whose interior contains  $\bar{x}$  but no point of  $\mathbb{Z}^p \times \mathbb{R}^{n-p}$  is contained in an inclusion *maximal* such set.
- How can we construct a (maximal) convex set  $C$  whose interior contains  $\bar{x}$  but no point of  $\mathbb{Z}^p \times \mathbb{R}^{n-p}$ ?

In the space  $\mathbb{R}^p$ , construct a  $\mathbb{Z}^p$ -free (maximal) convex set  $K$  whose interior contains the projection of  $\bar{x}$ .

The *cylinder*  $C = K \times \mathbb{R}^{n-p}$  is a (maximal) convex set whose interior contains  $\bar{x}$  but no point of  $\mathbb{Z}^p \times \mathbb{R}^{n-p}$ .

# EXAMPLE: Intersection Cuts from Split Disjunctions

Consider a split disjunction  $\pi x \leq \pi_0$  or  $\pi x \geq \pi_0 + 1$ ,  
where  $\pi \in \mathbb{Z}^P \times \{0\}^{n-p}$  and  $\pi_0 \in \mathbb{Z}$ , and  $\pi_0 < \pi \bar{x} < \pi_0 + 1$ .

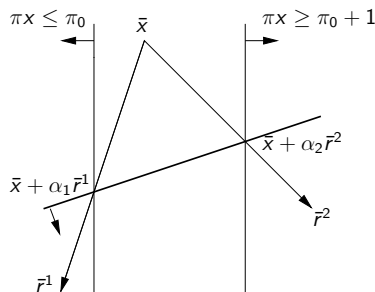
$K := \{x \in \mathbb{R}^P : \pi_0 \leq \sum_{j=1}^P \pi_j x_j \leq \pi_0 + 1\}$  is  $\mathbb{Z}^P$ -free and convex.

The set  $C := K \times \mathbb{R}^{n-p} = \{x \in \mathbb{R}^n : \pi_0 \leq \pi x \leq \pi_0 + 1\}$  is  $\mathbb{Z}^P \times \mathbb{R}^{n-p}$ -free.

Define  $\epsilon := \pi \bar{x} - \pi_0$ . We have  
 $0 < \epsilon < 1$ . For  $j \in N$ , define:

$$\alpha_j := \begin{cases} -\frac{\epsilon}{\pi \bar{r}^j} & \text{if } \pi \bar{r}^j < 0, \\ \frac{1-\epsilon}{\pi \bar{r}^j} & \text{if } \pi \bar{r}^j > 0, \\ +\infty & \text{otherwise,} \end{cases}$$

where  $\bar{r}^j$  are the rays of  $P(B)$ .



Intersection cut  $\sum_{j \in N} \frac{x_j}{\alpha_j} \geq 1$ .

## Gomory Mixed Integer Cuts from the Tableau

Let  $x_i, i \in B$  be a basic integer variable, and suppose  $\bar{x}_i = \bar{b}_i$  is fractional. We define  $\pi_0 := \lfloor \bar{x}_i \rfloor$ , and for  $j = 1, \dots, p$ ,

$$\pi_j := \begin{cases} \lfloor \bar{a}_{ij} \rfloor & \text{if } j \in N \text{ and } f_j \leq f_0, \\ \lceil \bar{a}_{ij} \rceil & \text{if } j \in N \text{ and } f_j > f_0, \\ 1 & \text{if } j = i, \\ 0 & \text{otherwise.} \end{cases}$$

For  $j = p + 1, \dots, n$ , we define  $\pi_j := 0$ .

Next we derive the intersection cut from the split disjunction  $\pi x \leq \pi_0$  or  $\pi x \geq \pi_0 + 1$  as shown in the previous slide. We need to compute  $\alpha_j, j \in N$  using our formula:

$$\alpha_j := \begin{cases} -\frac{\epsilon}{\pi \bar{r}^j} & \text{if } \pi \bar{r}^j < 0, \\ \frac{1-\epsilon}{\pi \bar{r}^j} & \text{if } \pi \bar{r}^j > 0, \\ +\infty & \text{otherwise,} \end{cases}$$

where  $\bar{r}^j$  are the rays of  $P(B)$ .

# Gomory Mixed Integer Cuts from the Tableau

Let  $f_0 = \bar{b}_i - \lfloor \bar{b}_i \rfloor$  and  $f_j = \bar{a}_{ij} - \lfloor \bar{a}_{ij} \rfloor$ . We have

$$\epsilon = \pi \bar{x} - \pi_0 = \sum_{h \in B} \pi_h \bar{x}_h - \pi_0 = \bar{x}_i - \lfloor \bar{x}_i \rfloor = f_0.$$

Let  $j \in N$ . We have  $\pi \bar{r}^j = \pi_j \bar{r}_j^j + \sum_{i \in N} \pi_i \bar{r}_i^j$  since  $\bar{r}_h^j = 0$  for all  $h \in N \setminus \{j\}$  and  $\pi_h = 0$  for all  $h \in B \setminus \{i\}$ . Therefore

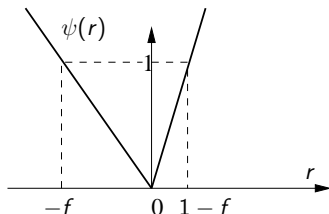
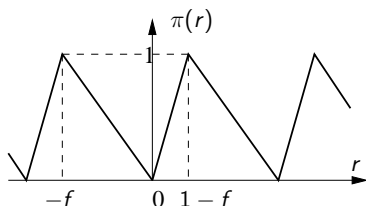
$$\pi \bar{r}^j = \begin{cases} \lfloor \bar{a}_{ij} \rfloor - \bar{a}_{ij} = -f_j & \text{if } 1 \leq j \leq p \text{ and } f_j \leq f_0, \\ \lceil \bar{a}_{ij} \rceil - \bar{a}_{ij} = 1 - f_j & \text{if } 1 \leq j \leq p \text{ and } f_j > f_0, \\ -\bar{a}_{ij} & \text{if } j \geq p + 1. \end{cases}$$

Now  $\alpha_j$  follows. Therefore the intersection cut associated with the split disjunction  $\pi x \leq \pi_0$  or  $\pi x \geq \pi_0 + 1$  is

$$\sum_{\substack{j \in N, j \leq p \\ f_j \leq f_0}} \frac{f_j}{f_0} x_j^+ + \sum_{\substack{j \in N, j \leq p \\ f_j > f_0}} \frac{1 - f_j}{1 - f_0} x_j^+ + \sum_{\substack{p+1 \leq j \leq n \\ \bar{a}_{ij} > 0}} \frac{\bar{a}_{ij}}{f_0} x_j^- - \sum_{\substack{p+1 \leq j \leq n \\ \bar{a}_{ij} < 0}} \frac{\bar{a}_{ij}}{1 - f_0} x_j \geq 1.$$

This is the **GMI cut**, since  $f_i = 0$  for  $i \in B$ .

# Gomory Functions



The Gomory formula looks complicated, and it may help to think of it as an inequality of the form

$$\sum_{j=1}^p \pi(\bar{a}_{ij})x_j + \sum_{j=p+1}^n \psi(\bar{a}_{ij})x_j \geq 1$$

where the functions  $\pi$  and  $\psi$ , are

$$\pi(a) := \min\left\{\frac{f}{f_0}, \frac{1-f}{1-f_0}\right\} \quad \text{and} \quad \psi(a) := \max\left\{\frac{a}{f_0}, \frac{-a}{1-f_0}\right\}.$$

with  $f = a - \lfloor a \rfloor$ .

## Maximal lattice-free convex sets

As observed earlier, the best possible intersection cuts are the ones defined by full-dimensional *maximal*  $(\mathbb{Z}^p \times \mathbb{R}^{n-p})$ -free convex sets in  $\mathbb{R}^n$ .

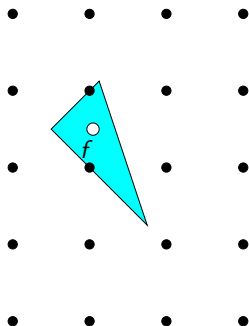
**LEMMA** Let  $C$  be a full-dimensional maximal  $(\mathbb{Z}^p \times \mathbb{R}^{n-p})$ -free convex set and let  $K$  be its projection onto  $\mathbb{R}^p$ . Then  $K$  is a full-dimensional maximal  $\mathbb{Z}^p$ -free convex set and  $C = K \times \mathbb{R}^{n-p}$ .

The next task is to understand the structure of full-dimensional maximal  $\mathbb{Z}^p$ -free convex sets.



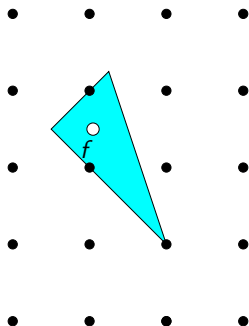
# Maximal Lattice-Free Convex Set

- ▶ Lattice-free convex set contains no integral point in its interior



# Maximal Lattice-Free Convex Set

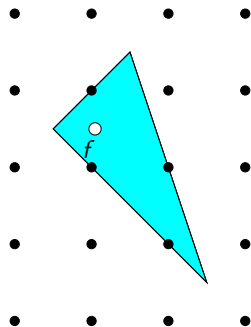
- ▶ **Lattice-free convex set** contains no integral point in its interior



**Maximal:** each edge contains an integral point in its relative interior.

# Maximal Lattice-Free Convex Set

- ▶ Lattice-free convex set contains no integral point in its interior



**Maximal:** each edge contains an integral point in its relative interior.

**In the plane:** it is a strip, a triangle or a quadrilateral.

## Theorem [Lovász 1989]

A maximal lattice-free convex set in the plane  $(x_1, x_2)$  is one of the following:

- i) Irrational line  $ax_1 + bx_2 = c$  with  $a/b$  irrational;
- ii) A strip  $c \leq ax_1 + bx_2 \leq c + 1$   
with  $a, b$  coprime integers,  $c$  integer;
- iii) A triangle with an integral point in the relative interior of each edge;
- iv) A quadrilateral containing exactly four integral points, one in the relative interior of each edge; The four integral points are vertices of a parallelogram of area 1.

# Lovász' Theorem

## THEOREM

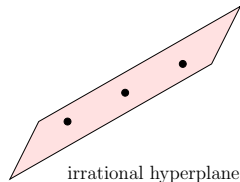
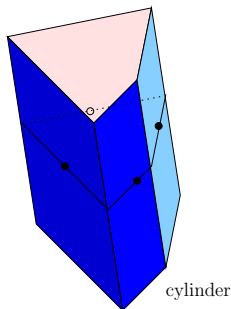
A set  $K \subset \mathbb{R}^p$  is a full-dimensional maximal  $\mathbb{Z}^p$ -free convex set if and only if

$K$  is a polyhedron of the form  $K = P + L$

where  $P$  is a polytope,  $L$  is a rational linear space,

$\dim(P) + \dim(L) = p$ ,

$K$  does not contain any point of  $\mathbb{Z}^p$  in its interior and there is a point of  $\mathbb{Z}^p$  in the relative interior of each facet of  $K$ .



# Proof of Lovász' Theorem

Let  $K \subset \mathbb{R}^p$  be a maximal  $\mathbb{Z}^p$ -free convex set.

We prove the theorem under the assumption that  $K$  is a **bounded set**. We need to show that  $K$  is a **polytope** and that **each of its facets has an integer point in its relative interior**.

Since we assume  $K$  bounded, there exist  $l, u$  in  $\mathbb{Z}^p$  such that  $K$  is contained in the box  $B = \{x \in \mathbb{R}^p : l_i \leq x_i \leq u_i, i = 1 \dots p\}$ .

Since  $K$  is a convex set, for each  $y \in B \cap \mathbb{Z}^p$ , there exists an half-space  $\{x \in \mathbb{R}^p : a_y x \leq b_y\}$  containing  $K$  such that  $a_y y = b_y$  (separation theorem for convex sets).

Since  $B$  is a bounded set,  $B \cap \mathbb{Z}^p$  is a finite set. Therefore  $P = \{x \in \mathbb{R}^p : l_i \leq x_i \leq u_i, i = 1 \dots p, a_y x \leq b_y, y \in B \cap \mathbb{Z}^p\}$  is a polytope.

By construction  $P$  is  $\mathbb{Z}^p$ -free and  $K \subseteq P$ .

Therefore  $K = P$  by maximality of  $K$ .

## Proof of Lovász' Theorem

We now show that each facet of  $K$  contains an integer point in its relative interior.

Suppose, by contradiction, that facet  $F_t$  of  $K$  does not contain a point of  $\mathbb{Z}^P$  in its relative interior.

Let  $a_t x \leq b_t$  be the inequality defining  $F_t$ .

Given  $\varepsilon > 0$ , let  $K'$  be the polyhedron defined by the same inequalities that define  $K$  except the inequality  $\alpha_t x \leq \beta_t$  that has been substituted with the inequality  $\alpha_t x \leq \beta_t + \varepsilon$ .

Since the recession cones of  $K$  and  $K'$  coincide,  $K'$  is a polytope. Since  $K$  is a maximal  $\mathbb{Z}^P$ -free convex set and  $K \subset K'$ ,  $K'$  contains points of  $\mathbb{Z}^P$  in its interior.

Since  $K'$  is a polytope, the number of points in  $K' \cap \mathbb{Z}^P$  is finite. Hence there exists one such point minimizing  $\alpha_t x$ , say  $z$ .

Let  $K''$  be the polytope defined by the same inequalities that define  $K$  except the inequality  $\alpha_t x \leq \beta_t$  that has been substituted with the inequality  $\alpha_t x \leq \alpha_t z$ .

By construction,  $K''$  does not contain any point of  $\mathbb{Z}^P$  in its interior and properly contains  $K$ , contradicting the maximality of  $K$ .

# Bound on the Number of Facets of Maximal $\mathbb{Z}^p$ -Free Polyhedra

Doignon 1973, Bell 1977 and Scarf 1977 show the following.

**THEOREM** Any full-dimensional maximal lattice-free convex set  $K \subseteq \mathbb{R}^p$  has at most  $2^p$  facets.

**PROOF** By Lovász' theorem, each facet  $F$  contains an integral point  $x^F$  in its relative interior. If there are more than  $2^p$  facets, then two integral points  $x^F$  and  $x^{F'}$  must be congruent modulo 2. Now their middle point  $\frac{1}{2}(x^F + x^{F'})$  is integral and it is in the interior of  $K$ , contradicting the fact that  $K$  is lattice-free.  $\square$



# Exercises

In "Courses Material" on the webpage

<http://eventos.cmm.uchile.cl/discretas2016/>

do the following exercises in Course Notes "Cutting planes in integer programming"

Exercise 3.1

Exercise 3.6

Exercise 3.8

Optional: Exercise 3.9